

VyZX

Formal Verification of a Graphical Language

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The ZX Calculus...

... is a graphical language for reasoning about quantum systems

ZX diagrams are open graphs consisting of green “Z” or red “X” “spiders” and “connections” between them



The Z and X spider

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ZX diagrams are **open graphs** consisting of **green** “Z” or **red** “X” “spiders” and “connections” between them

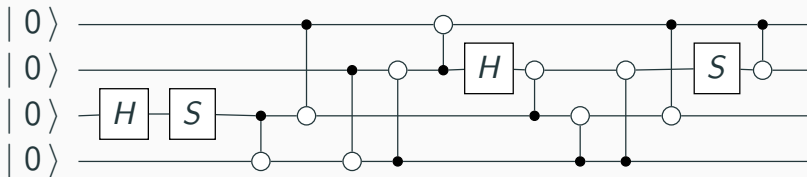


The Z and X spider

Used for **compilation, simulation, error correction** & more

Key benefit: **Diagrammatic rewrites** complete & more comprehensible than circuits or matrices

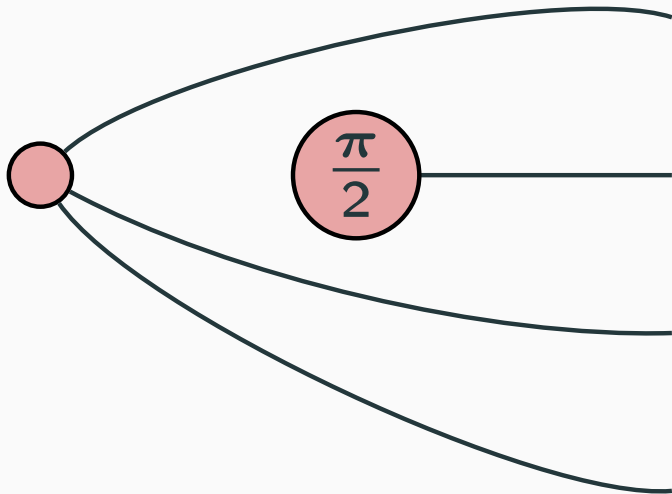
Example: Entanglement (Van de Wetering)



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$$[1 \ 0 \ 0 \ 0 \ i \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0 \ 0 \ i]^T$$

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- $|+\rangle, |-\rangle$ represents the X basis states (transform with hadamard (H))

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- For example, $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle, \alpha, \beta \in \mathbb{C}$.
- The bra notation $\langle\psi|$ is equivalent to $|\psi\rangle^T$.
- When measuring a qubit, it **collapses** to either $|0\rangle$ or $|1\rangle$.
- The probability of measuring $|0\rangle$ is $|\alpha|^2$.
- The probability of measuring $|1\rangle$ is $|\beta|^2$.
- The sum of probabilities is always 1: $|\alpha|^2 + |\beta|^2 = 1$.

Quantum Circuit Model

- Quantum operations are represented as gates in a circuit model.
- Each gate acts on qubits, changing their states.
- Common gates include Hadamard (H), Pauli-X (X), Pauli-Y (Y), and Pauli-Z (Z).
- Quantum circuits are read from left to right, and gates are applied in sequence.

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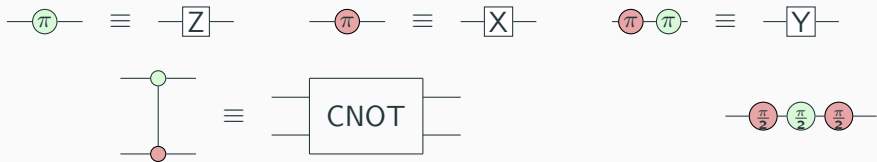
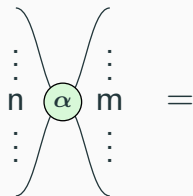
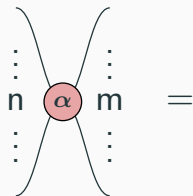


Diagram Semantics



$$\begin{array}{l} \begin{array}{c} \overset{n}{|0\rangle} \cdots \overset{n}{|0\rangle} \\ \vdots \\ \underset{n}{|1\rangle} \cdots \underset{n}{|1\rangle} \end{array} \mapsto \begin{array}{c} \overset{m}{|0\rangle} \cdots \overset{m}{|0\rangle} \\ \vdots \\ \underset{m}{e^{i\alpha} |1\rangle} \cdots \underset{m}{|1\rangle} \end{array} \end{array}$$



$$\begin{array}{l} \begin{array}{c} \overset{n}{|-\rangle} \cdots \overset{n}{|-\rangle} \\ \vdots \\ \underset{n}{|+\rangle} \cdots \underset{n}{|+\rangle} \end{array} \mapsto \begin{array}{c} \overset{m}{|-\rangle} \cdots \overset{m}{|-\rangle} \\ \vdots \\ \underset{m}{e^{i\alpha} |+\rangle} \cdots \underset{m}{|+\rangle} \end{array} \end{array}$$

Diagram Semantics

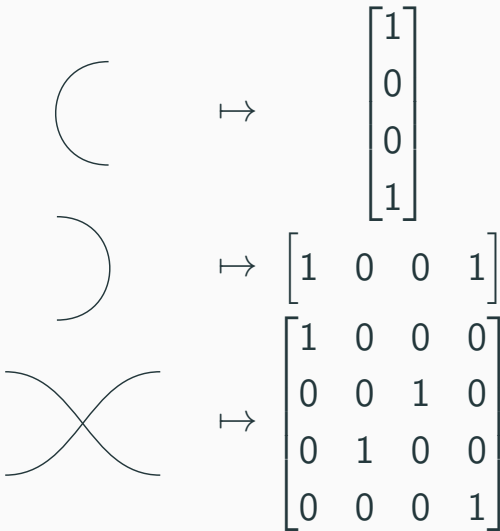
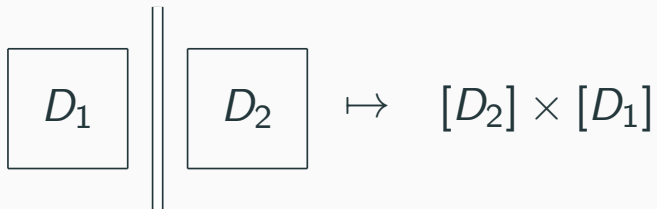
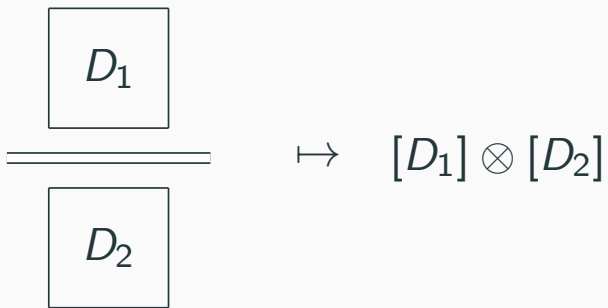


Diagram Semantics



A diagrammatic equation showing the semantics of the Cartesian product. On the left, two square boxes labeled D_1 and D_2 are placed side-by-side, separated by a vertical double line. An arrow points to the right, where the expression $[D_2] \times [D_1]$ is written.

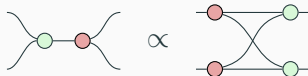
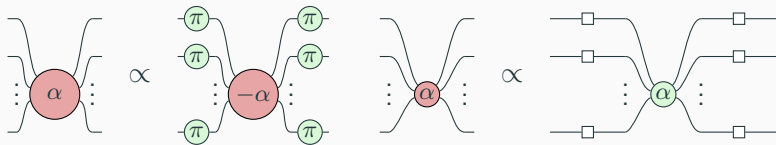
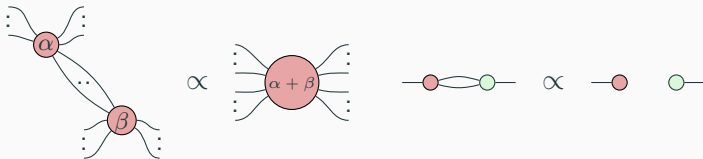
$$\boxed{D_1} \parallel \boxed{D_2} \mapsto [D_2] \times [D_1]$$



A diagrammatic equation showing the semantics of the tensor product. On the left, a square box labeled D_1 is positioned above a square box labeled D_2 , with a horizontal double line between them. An arrow points to the right, where the expression $[D_1] \otimes [D_2]$ is written.

$$\boxed{D_1} \overline{\hspace{2cm}} \boxed{D_2} \mapsto [D_1] \otimes [D_2]$$

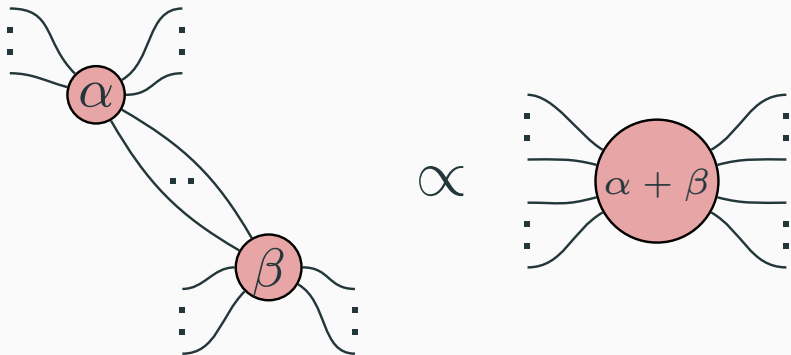
Rewriting Diagrams



Spider Fusion, The Hopf Rule, the Bi-pi Rule (Pi-Copy), Bi-hadamard Rule, the Bialgebra rule, and the identity rule

Spider Fusion

Spider fusion allows us to merge same colored spiders as long as they have one connection, adding their rotations together and connecting any inputs or outputs from the initial two spiders to the final spider.



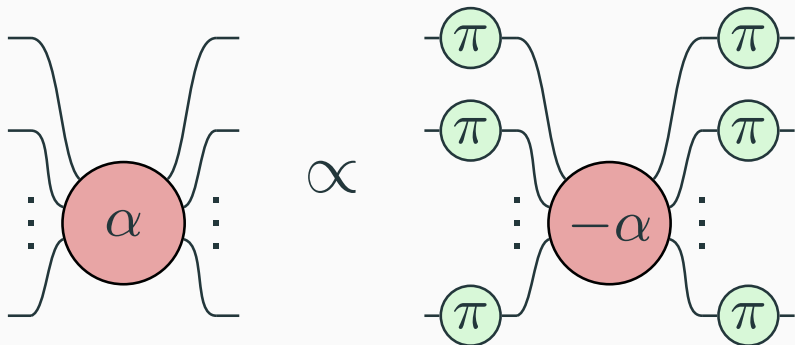
Hopf Rule

The hopf rule allows us to disconnect connections between opposite color spiders that come in pairs of two. A consequence of this with spider fusion is that n connections between opposite color spiders can always be considered to be equivalent to $n \bmod 2$ connections between those spiders.



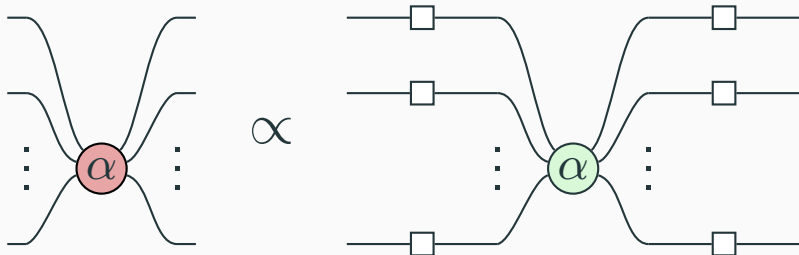
Bi-Pi Rule

The Bi-Pi rule allows us to add spiders of opposite colors with rotation π to every input and output of a spider and flip the phase.



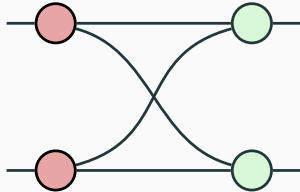
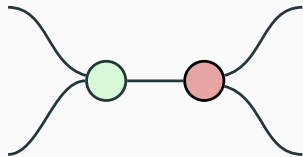
Bi-Hadamard Rule

The Bi-Hadamard rule allows us to add H-boxes to every input and output to flip the color of the spider within



Bialgebra rule

The bialgebra rule is unique in that it is one of the few rules that can introduce or remove swaps.



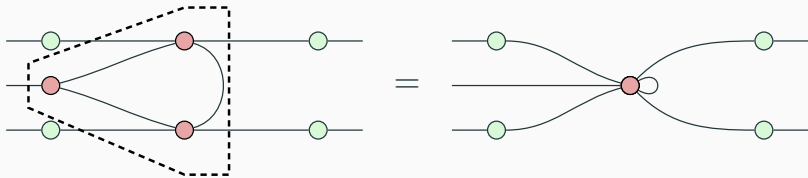
Identity Removal

The identity removal rules allow us to remove spiders with $k2\pi$ rotations in general.



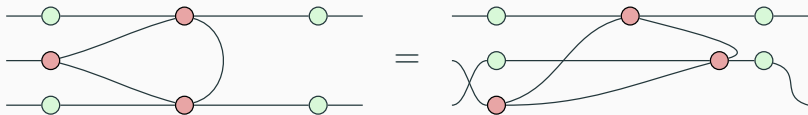
Rewriting Diagrams

Fusion can be used between the three connected X spiders here to simplify our diagram.

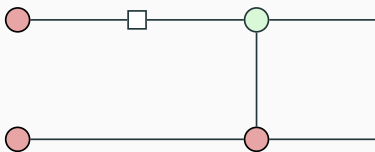
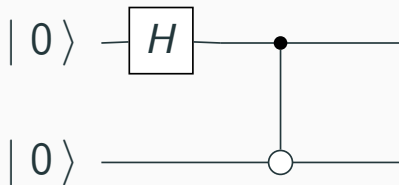


Only Connectivity Matters

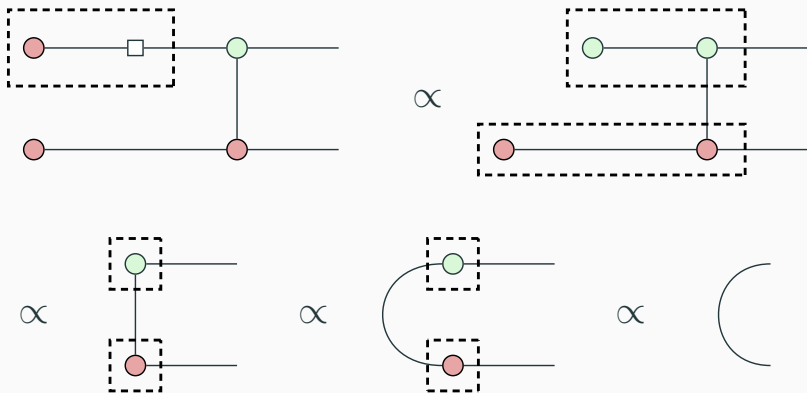
We can **freely move spiders around**, as long as their connections and in/outputs remain the same



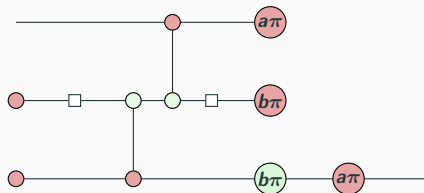
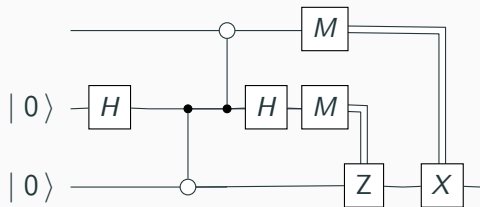
Example: Preparing a Bell Pair



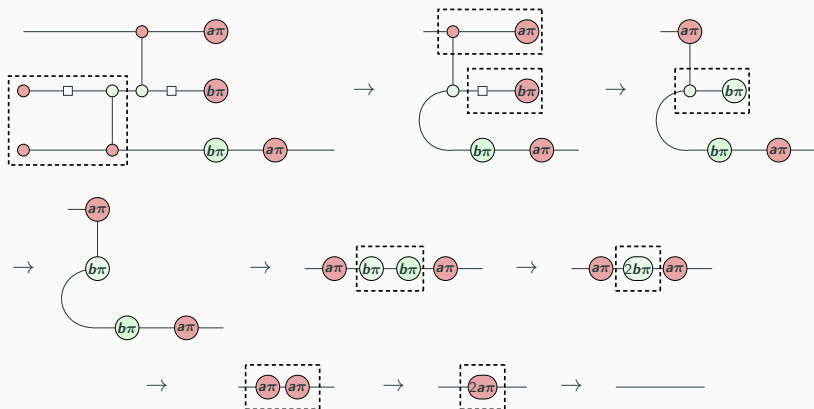
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Example: Teleportation (Van de Wetering)



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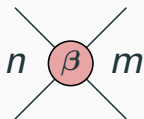
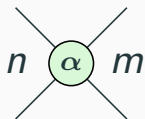
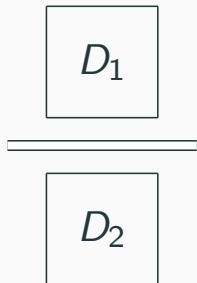
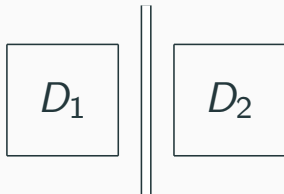
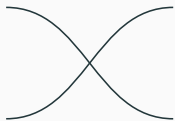
Unverified / Fully axiomatic:

- Quantomatic (<https://quantomatic.github.io/>)
- ZX Calculator (zx.cduck.me)
- Chyp

“Verify” by property testing:

- PyZX (<https://github.com/Quantomatic/pyzx>)

ZX Diagrams as string diagrams



Inductive ZX Diagrams

To define our ZX diagrams, we take these string diagram constructions and **add Z and X spiders**.

$$\begin{array}{c} \frac{\text{in out} : \mathbb{N} \quad \alpha : \mathbb{R}}{\text{Z_Spider in out } \alpha : \text{ZX in out}} \quad \frac{\text{in out} : \mathbb{N} \quad \alpha : \mathbb{R}}{\text{X_Spider in out } \alpha : \text{ZX in out}} \\ \\ \text{Cap} : \text{ZX } 0 \ 2 \quad \text{Cup} : \text{ZX } 2 \ 0 \quad \text{Swap} : \text{ZX } 2 \ 2 \quad \text{Empty} : \text{ZX } 0 \ 0 \\ \\ \frac{\text{zx1} : \text{ZX in mid} \quad \text{zx2} : \text{ZX mid out}}{\text{Compose zx1 zx2} : \text{ZX in out}} \quad \frac{}{\text{Wire} : \text{ZX } 1 \ 1} \\ \\ \frac{\text{zx1} : \text{ZX in1 out1} \quad \text{zx2} : \text{ZX in2 out2}}{\text{Stack zx1 zx2} : \text{ZX (in1 + in2) (out1 + out2)}} \quad \frac{}{\text{Box} : \text{ZX } 1 \ 1} \end{array}$$

To verify transformations on diagrams, we introduce a system of semantics. Our semantics system will rely on QuantumLib.

$$\begin{aligned} \text{Z_Spider } n \ m \ \alpha &\mapsto \begin{bmatrix} 1 & \cdots & 0 \\ \vdots & \ddots & 0 \\ 0 & 0 & e^{i\alpha} \end{bmatrix} \\ \text{X_Spider } n \ m \ \alpha &\mapsto H^{\otimes m} \times \begin{bmatrix} 1 & \cdots & 0 \\ \vdots & \ddots & 0 \\ 0 & 0 & e^{i\alpha} \end{bmatrix} \times H^{\otimes n} \end{aligned}$$

More Semantics

$$\text{Cap} \mapsto \begin{bmatrix} 1 & 0 & 0 & 1 \end{bmatrix}$$

$$\text{Cup} \mapsto \begin{bmatrix} 1 & 0 & 0 & 1 \end{bmatrix}^T$$

$$\text{Swap} \mapsto \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\text{Empty} \mapsto \begin{bmatrix} 1 \end{bmatrix}$$

$$\text{Wire} \mapsto I_{2 \times 2}$$

$$\text{Box} \mapsto H$$

Compose $zx1 \ zx2 \mapsto \text{semantics}(zx2) \times \text{semantics}(zx1)$

Stack $zx1 \ zx2 \mapsto \text{semantics}(zx1) \otimes \text{semantics}(zx2)$

Equivalence in ZX is up to constant factor

We define **proportionality** and use symbol \propto :

$$\exists c \neq 0 : \text{semantics}(zx1) = c * \text{semantics}(zx2) \implies zx1 \propto zx2$$

Allows Coq's **rewriting capabilities** in our proofs about diagrams

Why semantics?

- Smaller TCB
- Interoperability

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We can ingest quantum circuits using sqir

Circuit structure is very different

Convert circuit components

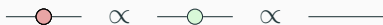
Prove equivalence through ground truth

Why semantics?

- Smaller TCB
- Interoperability
 - We can ingest quantum circuits using sqir
 - Circuit structure is very different
 - Convert circuit components
 - Prove equivalence through ground truth
- $VyZX$ can calculate results

Three Proof Strategies

1. Proof through semantics

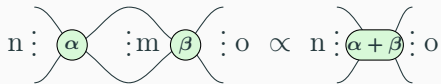


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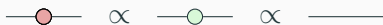


2. Inductive proof



Three Proof Strategies

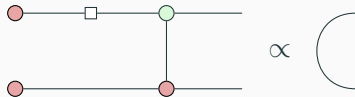
1. Proof through semantics



2. Inductive proof



3. Diagrammatic proof

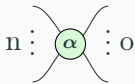


Three Proof Strategies

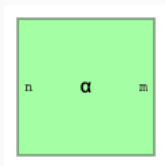
2. Inductive proof

$$n \vdots \alpha \quad \vdots m \quad \beta \quad \vdots o \quad \propto \quad n \vdots \alpha + \beta \quad \vdots o$$

Other representation



\cong

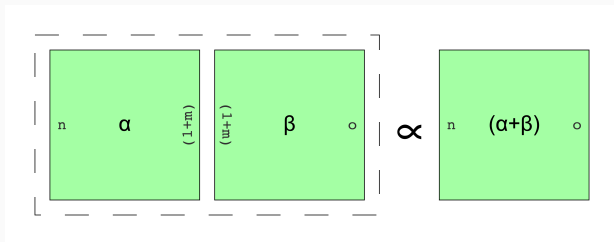


2. Inductive proof: Absolute Fusion



The diagram shows an equality between two configurations of wires and nodes. On the left, there are two nodes, α and β , each in a green circle. Node α has n input wires on the left and m output wires on the right. Node β has m input wires on the left and o output wires on the right. The m output wires of α are connected to the m input wires of β . On the right, there is a single node $\alpha + \beta$ in a green oval, with n input wires on the left and o output wires on the right. The two configurations are separated by a symbol \propto .

The lemma:



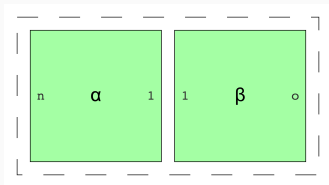
The diagram illustrates the lemma using rectangular nodes. On the left, two green rectangles are shown side-by-side, enclosed in a dashed box. The first rectangle has n input wires on the left, α in the center, and $(l+m)$ output wires on the right. The second rectangle has $(m+l)$ input wires on the left, β in the center, and o output wires on the right. On the right, a single green rectangle has n input wires on the left, $(\alpha+\beta)$ in the center, and o output wires on the right. The two configurations are separated by a symbol \propto .

Let's induct over m

2. Inductive proof: Absolute Fusion

$$n : \alpha : m : \beta : o \propto n : \alpha + \beta : o$$

Base case:

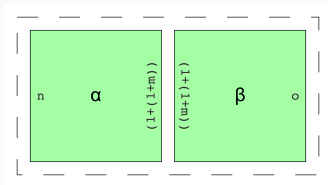


Solve with matrix semantics!

2. Inductive proof: Absolute Fusion

$$n \vdots \alpha \quad \vdots m \quad \beta \quad \vdots o \quad \propto \quad n \vdots \alpha + \beta \quad \vdots o$$

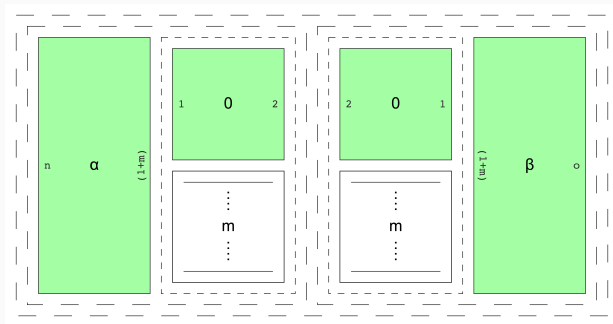
Inductive step:



Let's split out nodes to reduce size

2. Inductive proof: Absolute Fusion

$$n \text{ : } \alpha \text{ : } m \text{ : } \beta \text{ : } o \propto n \text{ : } \alpha + \beta \text{ : } o$$

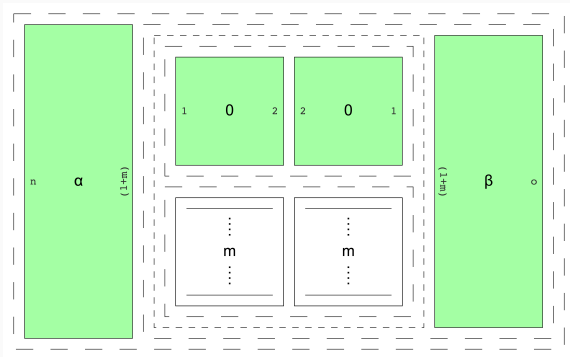


Idea: Fuse the small spiders. Problem: Association

2. Inductive proof: Absolute Fusion

$$n \text{ : } \alpha \text{ : } m \text{ : } \beta \text{ : } o \quad \propto \quad n \text{ : } \alpha + \beta \text{ : } o$$

Reassociated diagram

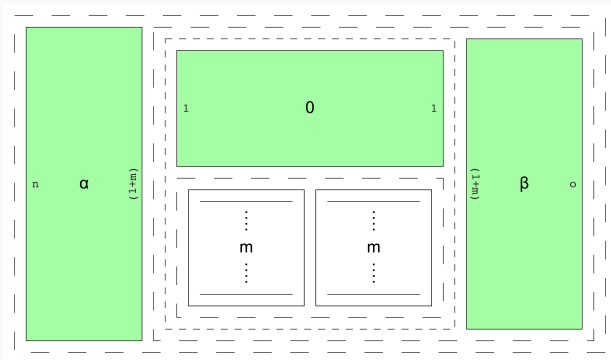


Let's fuse the small spiders

2. Inductive proof: Absolute Fusion

$$n \text{ : } \alpha \text{ : } m \text{ : } \beta \text{ : } o \quad \propto \quad n \text{ : } \alpha + \beta \text{ : } o$$

Upper spiders are fused

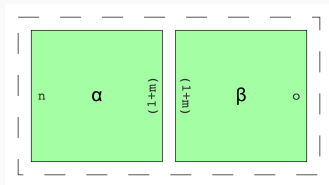


Apply identity rule, remove wires

2. Inductive proof: Absolute Fusion

$$n \vdots \alpha \vdots m \vdots \beta \vdots o \propto n \vdots \alpha + \beta \vdots o$$

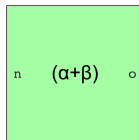
Diagram now corresponds to the IH



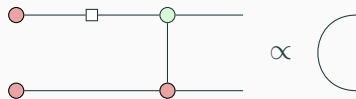
2. Inductive proof: Absolute Fusion

$$n \vdots \begin{array}{c} \curvearrowright \\ \alpha \\ \curvearrowleft \end{array} \vdots m \begin{array}{c} \curvearrowright \\ \beta \\ \curvearrowleft \end{array} \vdots o \propto n \vdots \begin{array}{c} \curvearrowright \\ \alpha + \beta \\ \curvearrowleft \end{array} \vdots o$$

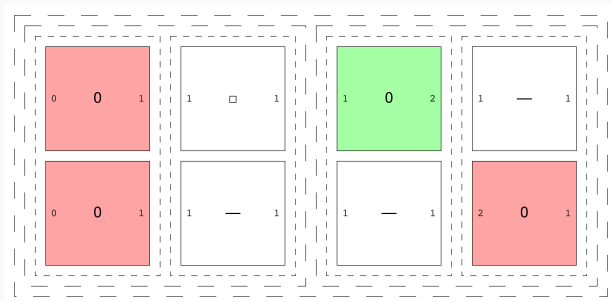
Qed.



3. Diagrammatic proof: Bell pair preparation

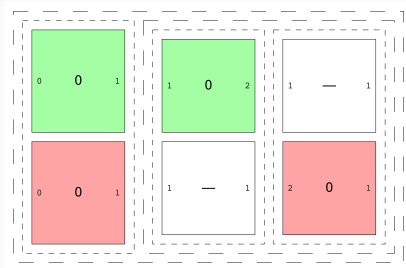
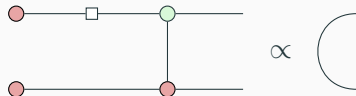


The lemma:



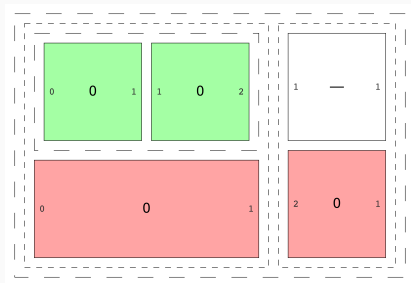
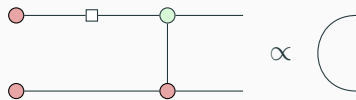
Let's first swap colors on the left

3. Diagrammatic proof: Bell pair preparation



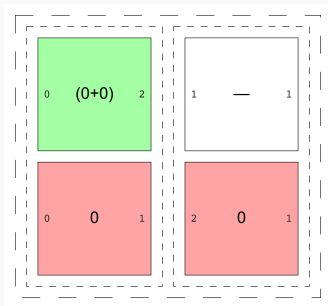
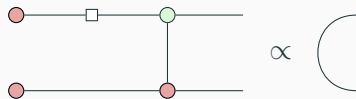
Let's reassociate!

3. Diagrammatic proof: Bell pair preparation



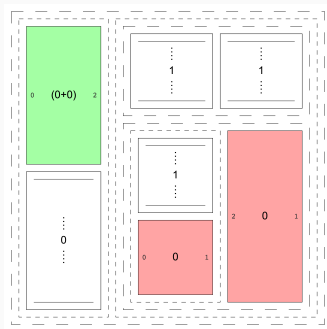
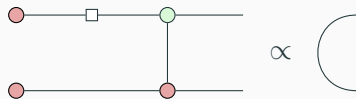
Let's fuse

3. Diagrammatic proof: Bell pair preparation



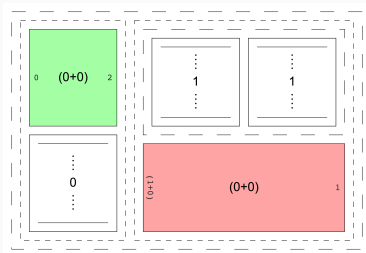
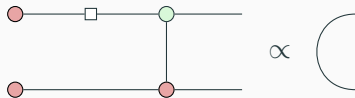
Separate the red node to fuse

3. Diagrammatic proof: Bell pair preparation

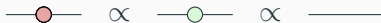


Let's fuse the red spiders

3. Diagrammatic proof: Bell pair preparation



Apply identity rule, remove wires



3. Diagrammatic proof: Bell pair preparation

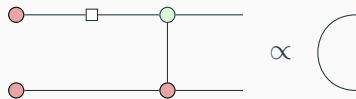
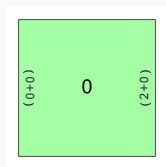
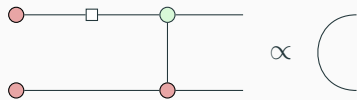


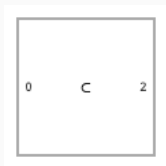
Diagram now corresponds to the a cap



3. Diagrammatic proof: Bell pair preparation



Qed.



Quantum circuit ingestion from RzQ gate set

$$\begin{array}{c} \frac{n < q : \mathbb{N}}{H\ n : \text{Circuit } q\ q} \\ \frac{n < q : \mathbb{N}}{X\ n : \text{Circuit } q\ q} \\ \frac{n < q : \mathbb{N} \quad \alpha : \mathbb{R}}{Rz(\alpha)\ n : \text{Circuit } q\ q} \\ \frac{c, n < q : \mathbb{N}}{CNOT\ c\ n : \text{Circuit } q\ q} \\ \frac{q : \mathbb{N} \quad c_1, c_2 : \text{Circuit } q\ q}{Compose\ c_1\ c_2 : \text{Circuit } q\ q} \end{array}$$

Quantum circuit ingestion from RzQ gate set

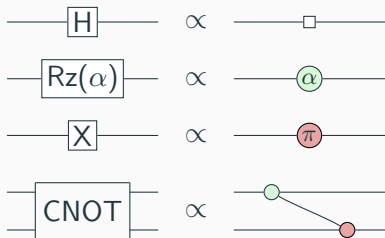
$$\frac{n < q : \mathbb{N}}{H \ n : \text{Circuit } q \ q}$$

$$\frac{n < q : \mathbb{N}}{X \ n : \text{Circuit } q \ q}$$

$$\frac{n < q : \mathbb{N} \quad \alpha : \mathbb{R}}{Rz(\alpha) \ n : \text{Circuit } q \ q}$$

$$\frac{c, n < q : \mathbb{N}}{CNOT \ c \ n : \text{Circuit } q \ q}$$

$$\frac{q : \mathbb{N} \quad c_1, c_2 : \text{Circuit } q \ q}{\text{Compose } c_1 \ c_2 : \text{Circuit } q \ q}$$



Quantum circuit ingestion from RzQ gate set

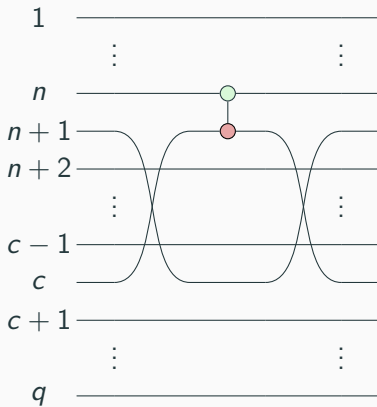
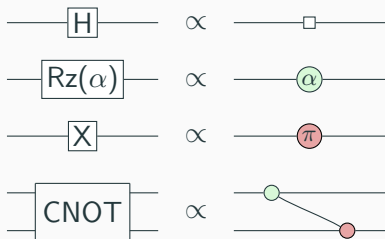
$\frac{n < q : \mathbb{N}}{H\ n : \text{Circuit } q\ q}$

 $\frac{n < q : \mathbb{N}}{X\ n : \text{Circuit } q\ q}$

 $\frac{n < q : \mathbb{N} \quad \alpha : \mathbb{R}}{Rz(\alpha)\ n : \text{Circuit } q\ q}$

$\frac{c, n < q : \mathbb{N}}{CNOT\ c\ n : \text{Circuit } q\ q}$

 $\frac{q : \mathbb{N} \quad c_1, c_2 : \text{Circuit } q\ q}{\text{Compose } c_1\ c_2 : \text{Circuit } q\ q}$



- Any ZX diagram can be expressed
- Multiple ways to encode
- Deal with associativity information
- Dimensionality issues

How can I verify my graphical language?

- Find underlying categorical structure
- Formally extend structure
- Translate into proof assistant
- Deal with resulting associativity issues

Future work

- Restore connection information
- Verify ZX-based compiler
- Prove ZX results

Summary

- Defined ZX diagrams inductively
- Inspired by string diagrams
- Multiple proof strategies
- ZX calculus is interesting!





Find VyZX on GitHub

<https://github.com/inQWIRE/VyZX>

arXiv

Coming soon...

References

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